Integration of quantum processes in cyber-physical systems

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Integration of quantum processes in cyber-physical systems

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- Emerged in the 19th century
 - railroads
 - telegraphic communications
- Popularized within computer science in the 60s by Dijkstra
- Nowadays is ubiquitous
 - web sites
 - mobile apps
- Concurrency:
 - is used to speed up processes
 - leads to notions of communication

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Usual Approach

Division of a task into independently small ones, which communicate with each other and are performed in an interleaved way

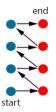


Figure: Example of a concurrent execution

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Concurrency: Problem

Memory access is a famous problem, due to:

- mutual exclusion
- deadlock
- starvation

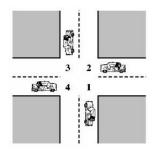


Figure: Cars on a crossing

Calculus of Communicating Systems

• One of the first algebras for reasoning about concurrency

• Developed by Robin Milner in 1989

• The central concept is that of a process, intuitively a black box

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CCS main constructors:

- . (action prefixing)
- + (choice operator; non-deterministic)
- | (parallel composition)

Example of a CCS process:

$$P \stackrel{def}{=} \overline{\texttt{coin}}.(\texttt{tea.P} + \texttt{coffee.P})$$

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Syntax fragment

P, Q ::= K | α .P | $\sum_{i \in I} P_i$ | P|Q

★ 3 → 3

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Operational Semantics

Set of rules that determines the behavior of a process

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The usual applications of process algebra are also important in quantum:

- quantum communication protocols
- concurrent quantum systems

• Developed by Mingsheng Ying et al in 2008

• No classical information involved

• Semantics in terms of labeled transition systems

• Extends CCS with super-operators

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qCCS: Underlying challenges

Restrictions from quantum mechanics:

• No cloning theorem

- Solved by syntactic restrictions
- Measurement destroys superpositions

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Syntax fragment

 $P,Q::=A(\tilde{x}) \mid nil \mid \varepsilon[X].P \mid c?x.P \mid c!x.P \mid P+Q \mid P \mid Q$

- $A(\tilde{x})$: constant processes
- $\varepsilon[X]$: action of a super-operator on a quantum state
- c?x: receive a quantum variable
- c!x: send a quantum variable

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Syntax fragment

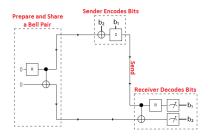
 $\mathsf{P},\mathsf{Q} ::= \mathsf{A}(\tilde{x}) \mid \mathtt{nil} \mid \varepsilon[X].\mathsf{P} \mid c?x.\mathsf{P} \mid c!x.\mathsf{P} \mid \mathsf{P}+\mathsf{Q} \mid \mathsf{P} ||\mathsf{Q}$

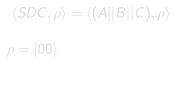
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The operational semantics is based on a configuration.

Configuration Pair formed by a process, P, and a state, ρ : $\langle P, \rho \rangle$

qCCS: Superdense Coding



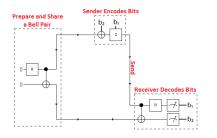


 $\begin{aligned} A &= d?x_1.(I[x_1].c!x_1.nil + Z[x_1].c!x_1.nil + X[x_1].c!x_1.nil + Z[x_1].X[x_1].c!x_1.nil) \\ B &= d?y_1.c?x_2.CNOT[x_2, y_1].H[x_2].M_{0,1}[x_2, y_1].nil \\ C &= c?x_0.c?y_0.H[x_0].CNOT[x_0, y_0].d!x_0.d!y_0.nil \end{aligned}$

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A B M A B M

qCCS: Superdense Coding



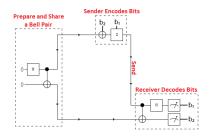
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 $\rho = |00 \rangle$

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Classically, the notion of time is essential to concurrent systems.

- is the basis of several synchronization mechanisms
- allows to reason about real timed systems
- and communication protocols
- All these aspects also apply to the quantum setting. Moreover, the notion of time also plays a key role in:
 - qubit coherence
 - gram architecture in the cloud

Our goal is to introduce a temporal version of qCCS.

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Examples

• Computer Science

 Control the time of calculations to see if it exceeded the coherence time of a qubit.

Physics

Photonic experiments, like quantum random number generator

Figure: Quantum Randon Number Generator

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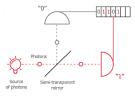


Figure: Quantum Randon Number Generator

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- should actions be considered instantaneous?
- should time domain be discrete?
- should time interfere with non-determinism?
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- temporal action that delays t units of time: $\sigma(t)$
- actions are considered to be performed instantaneously
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To do

- Syntax of Timed qCCS
- Operational semantics
- Notions of behavioral equivalence
- Answer the question
 - Do quantum laws interfere in the integration of time?

Thank you for your attention

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